Week 4 - Friday

#### Last time

- What did we talk about last time?
- Scheduling to minimize lateness
- Dijkstra's algorithm

#### **Questions?**



# **Big Oh**

- Order the following functions by rate of growth:
- $(\sqrt{2})^{\log n}$
- n<sup>2</sup>
- *n*!
- $(\log n)!$
- $\left(\frac{3}{2}\right)^n$
- n<sup>3</sup>
- $(\log n)^2$
- $\log(n!)$

- $2^{2^n}$
- $\log(\log n)$
- $n \cdot 2^n$
- $n^{\log(\log n)}$
- $\log n$
- **1**
- $2^{\log n}$
- $(\log n)^{\log n}$

- $4^{\log n}$
- (n+1)!
- $\sqrt{\log n}$
- *n*
- 2<sup>n</sup>
- $n \log n$



Determine the running time of various loops

• Example:

int counter = 0; for(int i = 0; i < n\*n; ++i) for(j = 1; j <= i; ++j) counter++;

# Graphs

- Know basic definitions of graphs
  - Nodes
  - Edges
  - Directed vs. undirected
  - Adjacency matrix vs. adjacency lists
  - Trees
  - Connected
  - Strongly connected

# Graph algorithms to know

- BFS
- DFS
- Determining bipartiteness
- Find connected components
- Find strongly connected components
- Topological sort

For all *n* ∈ Z, if *n*<sup>3</sup> + 5 is odd, then *n* is even.
Hint: Try a proof by contradiction.

• Prove that, for all integers  $n \ge 2$ ,

$$\sum_{i=1}^{n-1} i(i+1) = \frac{n(n-1)(n+1)}{3}$$

• **Hint:** Try a proof by induction.

• Prove that, for all integers  $n \ge 1$ ,

$$\frac{\sum_{i=1}^{n} 2i - 1}{\sum_{i=1}^{n} 2n + 2i - 1} = \frac{1}{3}$$

• **Hint:** Try a proof by induction.

- Prove that a graph with two or more nodes where each node has a degree of n/2 or higher must be connected.
- **Hint:** Try a proof by contradiction.

# Upcoming



#### Exam 1!

#### Reminders

- Finish Assignment 2
  - Due tonight before midnight
- Review chapters 1 through 3